Compiler Construction 2024-12-10



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Overview

Assembly in Practice

Summary

- · Assembly is the next step on the way to a working compiler
- x86 or ARM assembly can be quite complex and complicated
 - A relatively small subset should suffice for a minimalistic compiler
- We will mostly use GCC assembly for x86

- · Directives: Begin with a dot and indicate structural information
 - Example: .global main
- · Labels: End with a colon and denote special positions in the code
 - Example: main:
- · Instructions: Actual assembly code to be executed
 - Example: JMP main

Overview

Assembly in Practice

Summary

• Come up with an assembly version of the following for loop.

```
1 main: function void () = {
2     i: integer;
3     j: integer = 0;
4     for (i = 0; i < 42; i++) {
5          j++;
6     }
7 }</pre>
```

• Come up with an assembly version of the following for loop.

```
1
    main:
 2
         . . .
 3
        MOVL $0, %ebx # j
        MOVL $0, %eax # i
 4
 5
        JMP .for_check
 6
    .for_body:
 7
        ADDL $1. %ebx
 8
        ADDL $1, %eax
 9
    .for check:
        CMPL $41, %eax
10
11
        JLE .for_body
12
```

• Come up with an assembly version of the following switch statement.

```
1
    void foo (int i) {
         switch (i) {
 2
 3
              case 0:
 4
                   i++;
 5
                   break;
 6
              case 1:
 7
                   i - - :
 8
              case 2:
 9
                   i -= i:
10
         }
11
```

Group Exercise

• Come up with an assembly version of the following switch statement.

123456789 foo: # Assume i in %eax CMPL \$0, %eax JE .case_0 CMPL \$1. %eax JE .case_1 CMPL \$2, %eax JE .case 2 10 IMP end 11 12 .case_0: ADDL \$1, %eax 13 JMP .end 14 .case_1: 15 SUBL \$1, %eax 16 .case_2: 17 MOVL \$0. %eax 18 . end : 19

• Come up with an assembly version of alloca. Does it need any special handling?

```
void foo (void) {
       void* bar = alloca(128);
3
   }
```

1

2

• Come up with an assembly version of alloca. Does it need any special handling?

foo: 2 . . . 3 SUBO \$128, %rsp 4 MOVQ %rsp, bar 5

• alloca needs to be handled as a compiler built-in

1

Group Exercise

• What could go wrong in the following code snippet?

```
void foo (void) {
1
2
        void* bar;
 3
         . . .
4
        bar = alloca(1024);
5
         . . .
6
7
    int main (void) {
8
        for (int i = 0; i < 100000; i++) {
9
             foo();
10
         }
11
```

Group Exercise

• What could go wrong in the following code snippet?

```
void foo (void) {
 1
2
        void* bar;
 3
         . . .
        bar = alloca(1024);
4
 5
         . . .
6
7
    int main (void) {
8
        for (int i = 0; i < 100000; i++) {
9
             foo();
10
         }
11
```

· Inlining could lead to a stack overflow

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Assembly in Practice

Summary

- Assembly is CPU-specific
 - x86 and ARM support different instructions
- · There are different assembly dialects
 - Intel lists the source first, while AT&T/GCC lists the destination first
- · Different instruction suffixes and registers support different data sizes
 - · Several addressing modes and additional registers give more flexibility

[Thain, 2020] Thain, D. (2020). Introduction to Compilers and Language Design: Second Edition. http://compilerbook.org/.