Parsing (Part 1)

Compiler Construction

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Prof. Dr. Michael Kuhn
michael.kuhn@ovgu.de

Parallel Computing and I/O
Institute for Intelligent Cooperating Systems
Faculty of Computer Science
Otto von Guericke University Magdeburg
https://parcio.ovgu.de

Outline

- Terminal: Final symbol of the language (lower case letters)
- Non-terminal: Structure that can be transformed (upper case letters)
- Sentence: Sequence of terminals of the language
  - Sentential form: Sequence of terminals and non-terminals (Greek letters)
- Context-free grammar: List of rules making up the language
  1. $P \rightarrow E$ (P is the start symbol)
  2. $E \rightarrow X + Y$
  3. $X \rightarrow \text{int}$
  4. $Y \rightarrow \text{float}$

Group Exercise

- Perform a top-down derivation of $4 * \text{foo} + 1$ using grammar G_4.

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Sentential form	Rule
P	$P \rightarrow E$
E	$E \rightarrow E + T$
$E + T$	$E \rightarrow T$
$T + T$	$T \rightarrow T * F$
• $T * F + T$	$T \rightarrow F$
$F * F + T$	$F \rightarrow \text{int}$
$\text{int} * F + T$	$F \rightarrow \text{ident}$
$\text{int} * \text{ident} + T$	$T \rightarrow F$
$\text{int} * \text{ident} + F$	$F \rightarrow \text{int}$
$\text{int} * \text{ident} + \text{int}$	

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	Sentential form	Rule
	int $*$ ident $+$ int	$F \rightarrow$ int
	int $*$ ident $+ F$	$T \rightarrow F$
	int $*$ ident $+ T$	$F \rightarrow$ ident
	int $* F + T$	$F \rightarrow$ int
•	$F * F + T$	$T \rightarrow F$
	$T * F + T$	$T \rightarrow T * F$
	$T + T$	$E \rightarrow T$
	$E + T$	$E \rightarrow E + T$
	E	$P \rightarrow E$
	P	

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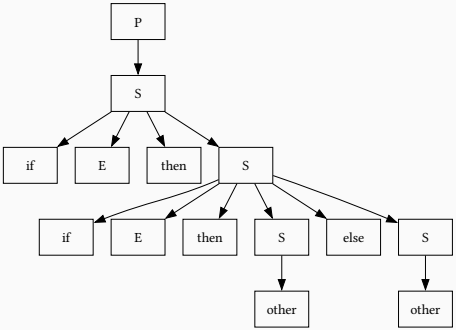
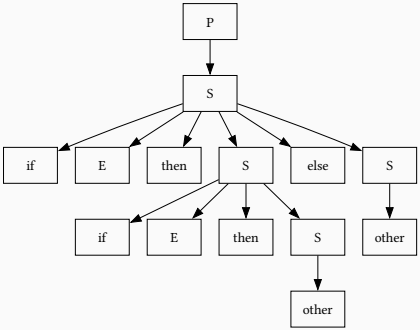
Sentential form	Rule
int $*$ ident $+$ int	$F \rightarrow$ int
int $*$ ident $+ F$	$T \rightarrow F$
int $*$ ident $+ T$	$F \rightarrow$ ident
• int $* F + T$	$T \rightarrow F$
int $* T + T$	$E \rightarrow T$
int $* E + T$	$E \rightarrow E + T$
int $* E$	$P \rightarrow E$
int $* P$	$\lightning$

Group Exercise

- Write out two possible parse trees using grammar G_5 for the sentence:
  `if E then if E then other else other`





- Modify grammar G_5 to prevent the dangling-else problem.
  (Hint: Prevent the inner S from containing an if without an else.)

- Modify grammar G_5 to prevent the dangling-else problem.
  (Hint: Prevent the inner S from containing an if without an else.)
  - $P \rightarrow S$
  - $S \rightarrow$ if E then S
  - $S \rightarrow S'$
  - $S' \rightarrow$ if E then S' else S
  - $S' \rightarrow$ other

Outline

LL Grammars

- LL(1) grammars are a subset of context-free grammars
- They can be parsed considering only one non-terminal and the next token
  - Remove ambiguity
  - Eliminate left recursion
  - Eliminate common left prefixes

Group Exercise

- Eliminate left recursion for grammar G_{12}.

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  - $E \rightarrow \text{id } E'$
  - $E \rightarrow \text{integer } E'$
  - $E' \rightarrow +EE'$
  - $E' \rightarrow \epsilon$

Group Exercise

- Eliminate common left prefixes for grammar G_{12}.

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  - $S \rightarrow$ if E then SS'
  - $S' \rightarrow \epsilon$
  - $S' \rightarrow$ else S

- $P \rightarrow S$
- $P \rightarrow SP$
- $S \rightarrow \text{if } E \text{ then } SS'$
- $S \rightarrow \text{while } ES$
- $S \rightarrow \text{begin } P \text{ end}$
- $S \rightarrow \text{print } E$
- $S \rightarrow E$
- $S' \rightarrow \epsilon$
- $S' \rightarrow \text{else } S$
- $E \rightarrow \text{id } E'$
- $E \rightarrow \text{integer } E'$
- $E' \rightarrow +EE'$
- $E' \rightarrow \epsilon$

- Determine the FIRST sets of the previous grammar.

Group Exercise

- Determine the FIRST sets of the previous grammar.
  - $P = \{\text{if}, \text{while}, \text{begin}, \text{print}, \text{id}, \text{integer}\}$
  - $S = \{\text{if}, \text{while}, \text{begin}, \text{print}, \text{id}, \text{integer}\}$
  - $S' = \{\text{else}, \epsilon\}$
  - $E = \{\text{id}, \text{integer}\}$
  - $E' = \{+, \epsilon\}$

- Determine the FOLLOW sets of the previous grammar.

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  - $P = \{\$\}$
  - $S = \{\$, \text{if}, \text{while}, \text{begin}, \text{print}, \text{id}, \text{integer}, \text{else}\}$
  - $S' = \{\$, \text{if}, \text{while}, \text{begin}, \text{print}, \text{id}, \text{integer}, \text{else}\}$
  - $E = \{\text{then}, \text{if}, \text{while}, \text{begin}, \text{print}, \text{id}, \text{integer}, \$, \text{else}, +\}$
  - $E' = \{\text{then}, \text{if}, \text{while}, \text{begin}, \text{print}, \text{id}, \text{integer}, \$, \text{else}, +\}$

Outline

- Context-free grammars are more powerful than regular expressions
  - They are described using multiple rules that replace non-terminals
- LL(1) are a subset of context-free grammars without ambiguities
  - They also eliminate left recursion and common left prefixes

References

[Thain, 2020] Thain, D. (2020). ***Introduction to Compilers and Language Design: Second Edition.*** http://compilerbook.org/.